

Discrete Optimization with Ordering

Elena Fernández

Barcel-berg { Technical University of Catalonia
Visiting Institut für Informatik. Heidelberg

Justo Puerto

University of Seville

Antonio Rodríguez-Chía

University of Cádiz

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Discrete Optimization with ordering

Discrete optimization problems where feasible solutions are sequences of elements which are ordered with respect to a priority (hierarchy) function.

The cost of an element depends on its position on the sequence.

- Multiperiod problems
- Scheduling and sequencing problems
- ...

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- The simple ordering problem (SOP)
 - Ordered sequences
 - Some properties
 - The polyhedron of the SOP
 - The simple ordering problem with cardinality constraint
 - The simple ordering problem on an independence system
 - The ordered median spanning tree problem

Ordered sequences

- **Ground set:** $E = \{e_1, e_2, \dots, e_n\}$. $N = \{1, 2, \dots, n\}$.
- $c: E \rightarrow \mathbb{R}$ **order function** s.t. $c_1 := c(e_1) \geq c(e_2) \geq \dots \geq c(e_n) := c_n$
- **Feasible Solutions:** sequences with at most $p \leq n$ elements
 which are ordered wrt function c . $K = \{1, 2, \dots, p\}$.
 $[e_{j_1}, e_{j_2}, \dots, e_{j_r}]$, $r \leq p$, such that $j_i < j_{i+1}$



- E_K : multiset with p copies of each element $e \in E$.
 - $F \subseteq E_K$, $F = \{e_{j_1}^{k_1}, e_{j_2}^{k_2}, \dots, e_{j_r}^{k_r}\}$ with $k_1 \leq k_2 \leq \dots \leq k_r$
 - $F \subseteq E_K$, **ordered sequence** $\Leftrightarrow k_i < k_{i+1}$ and $j_i < j_{i+1}$, $i = 1, \dots, p-1$
- **Additive objective function:** The value of each element depends on its position in the sequence.

$$\begin{aligned}
 d: E_K &\longrightarrow \mathbb{R} \\
 e_j^k &\longrightarrow d_j^k \\
 F &\longrightarrow \sum_{e_j^k \in F} d_j^k
 \end{aligned}$$

Example

$$E = \{e_1, e_2, e_3, e_4\} (c_1 \geq c_2 \geq c_3 \geq c_4), p = 3.$$

$$d = \begin{pmatrix} 2 & 0 & 5 \\ 2 & 2 & 3 \\ 3 & 0 & 1 \\ 1 & 1 & 2 \end{pmatrix}$$

$F = \{e_2^1, e_2^2, e_3^2, e_1^3\}$ is not an ordered sequence

Ordered sequences:

$$F_1 = \{e_2^1, e_3^2\}, \quad F_2 = \{e_2^2\}, \quad F_3 = \{e_1^3\}$$

$$d(F_1) = 2, \quad d(F_2) = 2, \quad d(F_3) = 5$$

Some properties of ordered sequences

- $I = (E_K, \mathcal{F})$ is an Independence System,

where $\mathcal{F} = \{F \subseteq E_K : F \text{ is an ordered sequence}\}$.

$F \subseteq E_K$ ordered sequence $\Rightarrow S$ ordered sequence, for all $S \subseteq F$.

- $\ell(F) = \max\{|S| : S \subseteq F, S \text{ is an ordered sequence}\}$.

F is an ordered sequence $\Leftrightarrow \ell(F) = |F|$.

- $I = (E_K, \mathcal{F})$ is not a matroid

For a given F , it is possible to find maximal ordered sequences $S, T \subseteq F$ such that $\ell(S) \neq \ell(T)$,

The Simple Ordering Problem (SOP):

Given $E, c, p, d,$

to find an ordered sequence of maximum total weight with respect to $d.$

$$\begin{array}{ll} d(F^*) := \max & d(F) \\ \text{s.t.} & |F| \leq \ell(F), \quad \text{for all } F \subseteq E_K \end{array}$$

An optimal solution may have any number of elements in the range $[0, p].$

$$d = \begin{pmatrix} -2 & 0 & -5 \\ -2 & -2 & -3 \\ -3 & 0 & -1 \\ -1 & -1 & -2 \end{pmatrix}$$

$$d = \begin{pmatrix} 2 & 0 & 15 \\ 2 & 2 & 3 \\ 3 & 0 & 1 \\ 1 & 1 & 2 \end{pmatrix}$$

$$d = \begin{pmatrix} 2 & 0 & 5 \\ 2 & 2 & 3 \\ 3 & 0 & 1 \\ 1 & 1 & 2 \end{pmatrix}$$

The polyhedron of the SOP

$$x_{jk} = \begin{cases} 1 & \text{if element } e_j^k \text{ is selected} \\ 0 & \text{otherwise} \end{cases}$$

$$P_{\text{SOP}} = \text{conv} \{x \in \{0, 1\}^{n \times p} : x(F) \leq \ell(F), \text{ for all } F \subseteq E_K\}$$

$$\triangleright \sum_{j \in N} x_{jk} \leq 1, \quad k \in K$$

$$\triangleright \sum_{k \in K} x_{jk} \leq 1, \quad j \in N$$

$$\triangleright x_{jk} + \sum_{j' \leq j} x_{j'k'} \leq 1, \quad \forall j \in N, \quad \forall k, k' \in K, k' > k$$

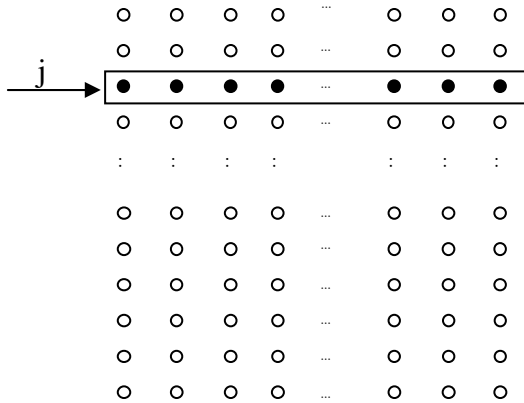


$$\sum_{j' \geq j} x_{j'k} + \sum_{j' \leq j} x_{j'k'} \leq 1, \quad \forall j \in N, \quad \forall k, k' \in K, k' > k$$

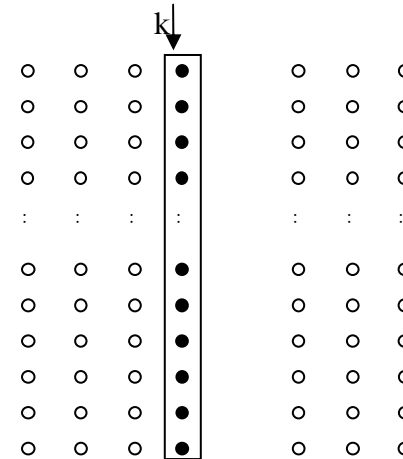


$$\sum_{j' \geq j} x_{j'k} + \sum_{k < h < k'} x_{jh} + \sum_{j' \leq j} x_{j'k'} \leq 1, \quad \forall j \in N, \quad \forall k, k' \in K, k' > k$$

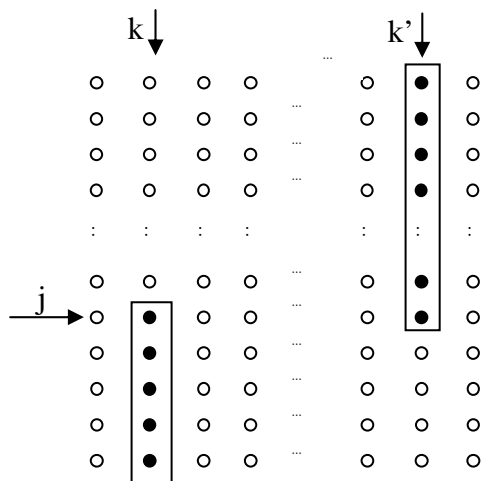
$$\sum_{k \in K} x_{jk} \leq 1$$



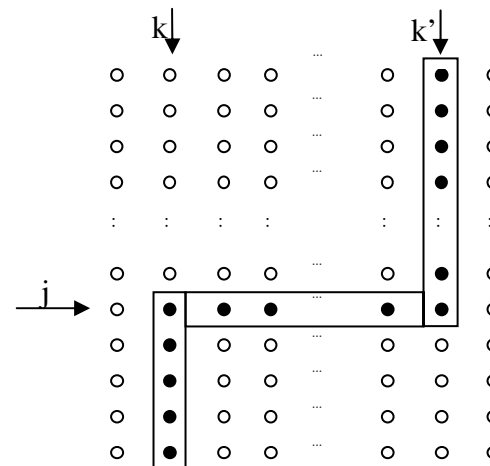
$$\sum_{j \in N} x_{jk} \leq 1$$



$$\sum_{j' \geq j} x_{j'k} + \sum_{j' < j} x_{j'k'} \leq 1$$



$$\sum_{j' \geq j} x_{j'k} + \sum_{k < h < k'} x_{jh} + \sum_{j' < j} x_{j'k'} \leq 1$$



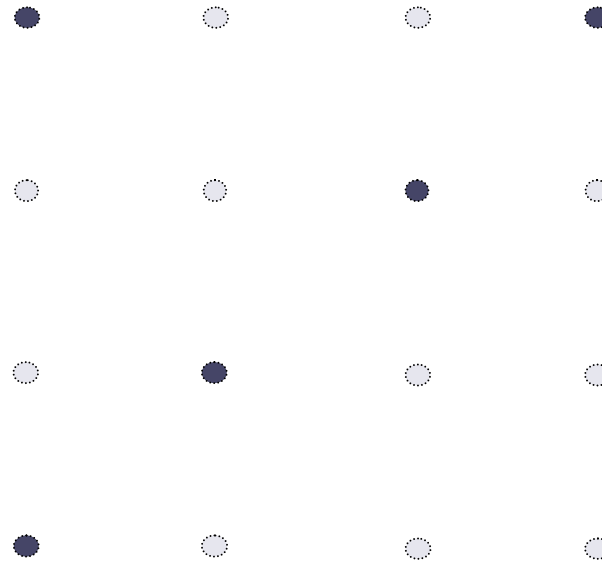
$$\sum_{j' \geq j} x_{j'k} + \sum_{k < h < k'} x_{jh} + \sum_{j' \leq j} x_{j'k'} \leq 1, \quad \forall j \in N, \quad \forall k, k' \in K, k' > k$$

$$\sum_{k' \leq k} x_{jk'} + \sum_{j < h < j'} x_{hk} + \sum_{k' > k} x_{j'k'} \leq 1 \quad \forall j, j' \in N, j' < j, \quad \forall k \in K$$

Not
enough

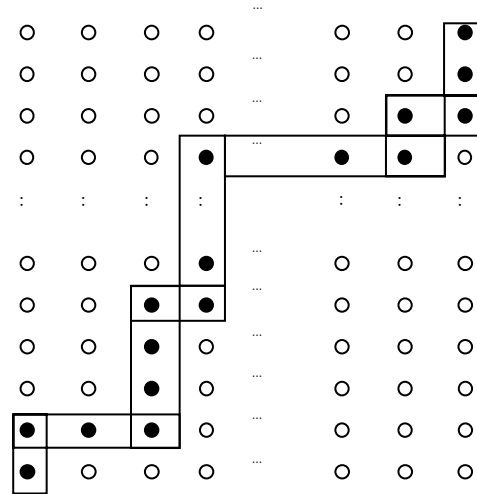
$$E = \{e_1, e_2, e_3, e_4\} \quad p=4,$$

$$x_{11} = x_{41} = x_{32} = x_{23} = x_{14} = 1/3$$



Staircase: $H \subseteq E_K$ s.t. $e_j^k, e_j^{k'} \in H$ and $j' \leq j$, then $k \leq k'$.

Maximal staircase: staircase not contained in any other



Staircase inequality

$$\sum_{e_j^k \in H} x_{jk} \leq 1$$



Valid inequality for the SOP that generalizes all the previous ones

H maximal staircase



undominated inequality

Theorem:

$$P_{\text{SOP}} = \{x \in [0, 1]^{n \times p} : x(H) \leq 1, \text{ for all maximal staircase } H \subseteq E_K\}$$

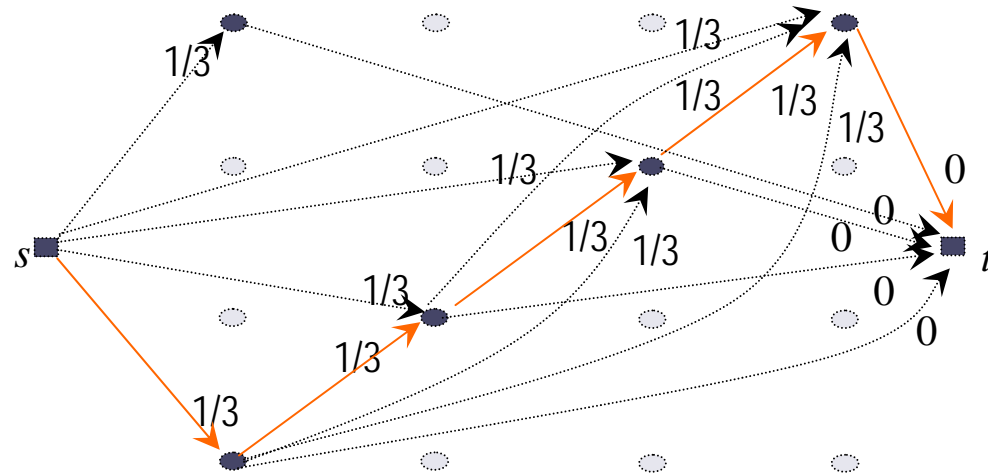
Separation of staircase inequalities

Proposition:

For a given $y \in \mathbb{R}^{n \times p}$ such that $0 \leq y_{jk} \leq 1$, for all $j \in N$, $k \in K$, the separation problem for staircase inequalities can be solved in polynomial time by finding an s - t -path of maximum cost in $N(y)$.

$$E = \{e_1, e_2, e_3, e_4\} \quad p=4,$$

$$x_{11} = x_{41} = x_{32} = x_{23} = x_{14} = 1/3$$



$$N(y) = (V_y \cup \{s, t\}, A(y))$$

- $V_y = \{v_j^k \in E_K : y_{jk} > 0\}$: support of y ;
- s and t : fictitious source and sink.

$A(y)$ contains the following arcs:

- One arc (s, v_j^k) of cost y_{jk} , for each node $v_j^k \in V_y$.
- One arc $(v_j^k, v_{j'}^{k'})$ of cost $y_{j'k'}$ for each pair $v_j^k, v_{j'}^{k'} \in V_y$, with $j \geq j'$ and $k \leq k'$.
- One arc (v_j^k, t) of cost zero, for each node $v_j^k \in V_y$.

Theorem:


The SOP can be solved in polynomial time.

The simple ordering problem with cardinality constraint (SOPC)

Given $E, c, p, d,$

to find an ordered sequence that contains exactly one element of each $k \in K$, of maximal total weight with respect to d .

$$\begin{aligned}
 d(F^*) &:= \max && d(F) \\
 &\text{s.t.} && |F| \leq \ell(F), && F \subseteq E_K. \\
 &&& |F|=p
 \end{aligned}$$

$ \begin{aligned} \sum_{e_j^k \in H} x_{jk} &\leq 1 && \forall H \text{ maximal stair} \\ \sum_{j \in N} x_{jk} &= 1, && k \in K \end{aligned} $		$ \begin{aligned} \sum_{e_j^k \in H} x_{jk} &\leq 1 && \forall H \text{ maximal stair} \\ \sum_{k \in K} \sum_{j \in N} x_{jk} &= p \end{aligned} $
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Theorem:

$$P_{\text{SOPC}} = \{x \in [0, 1]^{n \times p} : x(H) \leq 1, \text{ for all maximal staircase } H \subseteq E_K, \sum_{k \in K} \sum_{j \in N} x_{jk} = p \}.$$

Theorem:

The SOPC can be solved in polynomial time.

The simple ordering problem on an independence system (SOPI)

Given E, c, p, d + Independence System $J = (E, \mathcal{H})$

$I = (E_K, \mathcal{F})$ Independence System induced by order function c

$J_K = (E_K, \mathcal{H}_K)$ Independence System

$$\left\{ e_{j_1}^{k_1}, e_{j_2}^{k_2}, \dots, e_{j_r}^{k_r} \right\} \in \mathcal{H}_K \iff \left\{ e_{j_1}, e_{j_2}, \dots, e_{j_r} \right\} \in \mathcal{H}$$

to find an ordered sequence of \mathcal{F} which is an independent set of \mathcal{H}_K of maximum total weight with respect to d .

$$\mathcal{L}(F) = \max\{|S| : S \subseteq F, S \in \mathcal{F} \cap \mathcal{H}_K\}.$$

$$d(F^*) := \max_{\substack{d(F) \\ \text{s.t. } |F| \leq \mathcal{L}(F), \\ F \subseteq E_K.}}$$

$$|F|=p$$

$$P_{\text{SOPI}} = \text{conv} \{x \in \{0, 1\}^{n \times p} : x(F) \leq \mathcal{L}(F), F \subseteq E_K\}$$

$$P_{\text{SOPI C}} = \text{conv} \{x \in \{0, 1\}^{n \times p} : x(F) \leq \mathcal{L}(F), F \subseteq E_K, |F|=p\}$$

Mathematical Programming Formulation of SOPI(C)

$$\begin{array}{ll}
 \text{Max} & \sum_{k \in K} \sum_{j \in N} d_j^k x_{jk} \\
 \text{s.t.} & x(H) \leq 1 \quad \text{for all maximal staricase } H \subset E_K \quad (1) \\
 & x(F) \leq r(F) \quad \text{for all } F \subseteq E_K \quad (2) \\
 & x(F) = p \quad (3) \\
 & x_{jk} \in \{0,1\} \quad \text{for all } j \in N, k \in K
 \end{array}$$

$r(\cdot)$: rank function on J_K

$\{x \in \{0, 1\}^{n \times p} : x \text{ satisfies (1) and (2)}\}$

has fractional vertices

$E = \{e_1, e_2, e_3\}$ edges of K_3 , $p=3$ and \mathcal{H} given by forests in K_3

$$d = \begin{pmatrix} 1 & 0 & 1.5 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

$$x_{11} = x_{22} = x_{13} = x_{33} = 1/2$$

$$z^* = 2.25$$

Some properties of $P_{\text{SOPI}} = \text{conv} \{x \in \{0, 1\}^{n \times p} : x \text{ satisfies (1) and (2)}\}$

➤ Inequalities $x(F) \leq r(F)$ need not be facets of P_{SOPI}

$E = \{e_1, e_2, e_3\}$ edges of K_3 , $p=3$ and \mathcal{H} given by forests in K_3

$$F = \{e_2^1, e_2^2, e_3^3\}, \quad r(F) = 2$$

$$x_{11} + x_{22} + x_{33} \leq 2 \quad \text{is dominated by}$$

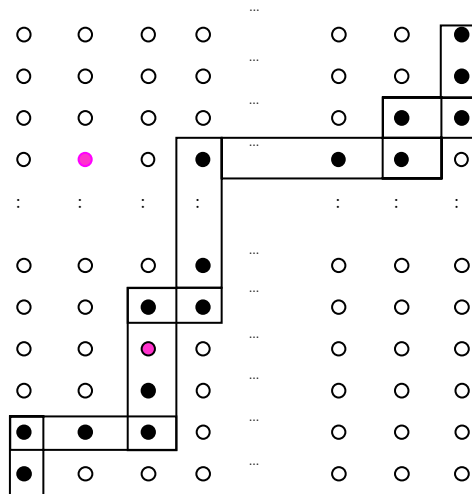
$$x_{11} + x_{22} + x_{32} + x_{33} \leq 2$$

More properties of $P_{\text{SOPI}} = \text{conv} \{x \in \{0, 1\}^{n \times p} : x \text{ satisfies (1) and (2)}\}$

Proposition: Let $H \subset E_K$ be a maximal staircase such that

for all $e_j^k \in E_K \setminus H \exists e_{j'}^{k'} \in H$ s.t. $\{e_j^k, e_{j'}^{k'}\} \in \mathcal{F} \cap \mathcal{H}_K$.

Then, $x(H) \leq 1$ is a facet of P_{SOPI} .



More properties of $P_{\text{SOPI}} = \text{conv} \{x \in \{0, 1\}^{n \times p} : x \text{ satisfies (1) and (2)}\}$

Proposition: Let $H \subset E_K$ be a maximal staircase such that
for all $e_j^k \in E_K \setminus H \exists e_{j'}^{k'} \in H$ s.t. $\{e_j^k, e_{j'}^{k'}\} \in \mathcal{F} \cap \mathcal{H}_K$.
Then, $x(H) \leq 1$ is a facet of P_{SOPI} .

Corollary: If $\{e_j, e_{j'}\} \in \mathcal{H}$ for all $j, j' \in N$, then for all maximal staircase H ,
 $x(H) \leq 1$ is a facet of P_{SOPI} .

Conditions hold:

- $J=(E, \mathcal{H})$: forests in graph (V, E)
- $J=(E, \mathcal{H})$: sets of l.i. columns of an $m \times n$ matrix A such that not two columns are l.d.
- $J=(E, \mathcal{H})$: independence system induced by a knapsack type constraint with coefficients vector $a \in \mathbb{R}^n$ and right hand side a_0 such that $a_j + a_{j'} \leq a_0$ for all $j, j', j \neq j'$.

Conditions do not hold:

- $J=(E, \mathcal{H})$: matchings in graph (V, E)

Ordered median objective function

Given	Ordered median of F
E, p, K $c: E \rightarrow \mathbb{R},$ $\lambda: \{1, 2, \dots, p\} \longrightarrow \mathbb{R}^+$ $F \subseteq E, F =p,$	$om(F) = \sum_{k \in K} \lambda_k c_{e_{\pi_F(k)}}$ $\pi_F: K \longrightarrow K \text{ t.q.}$ $c_{\pi_F(1)} \geq c_{\pi_F(2)} \geq \dots \geq c_{\pi_F(p)}.$

If	$\lambda=(1, 1 \dots, 1, \dots, 1)$	Sum	(median)
If	$\lambda=(1, 0 \dots, 0, \dots, 0)$	Maximum	(center)
If	$\lambda=(1, 1, \dots, 1, \dots, 0)$	k-center	
If	$\lambda=(1, \alpha, \dots, \alpha, \dots, \alpha)$	Cent dian	

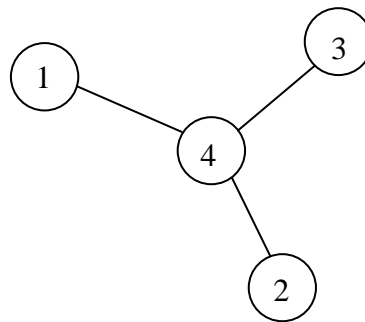
Nickel, S., and Puerto, J. (2005). Location Theory. A Unified approach. Springer.

The ordered median spanning tree problem

Given $G=(V, E)$, $c = (c_1, \dots, c_n)$, $c_1 \geq \dots \geq c_n$, $\lambda = (\lambda_1, \dots, \lambda_p) \geq 0$, the Ordered Median Spanning Tree Problem (OMSTP) consists of finding an ordered sequence $\{e^{k1}_{j_1}, \dots, e^{kp}_{j_p}\}$ such that $T = \{e_{j_1}, \dots, e_{j_p}\}$ is an spanning tree of G , that maximizes the value of $\sum_{k \in K} \lambda_k c_{j_k}$.

	2	3	4
1	4	3	5
2		6	8
3			7

$$\lambda_1=10, \lambda_2=1, \lambda_3=5$$



$$om(T) = 10 \times 8 + 1 \times 7 + 5 \times 5$$

$$c_{24} \geq c_{34} \geq c_{14}$$

The OMSTP

$$\begin{aligned} \text{Max} \quad & \sum_{k \in K} \sum_{j \in N} \lambda_k c_j x_{jk} \\ \text{s.t.} \quad & x(H) \leq 1 \quad \text{for all maximal staricase } H \subset E_K \\ & x(F) \leq r(F) \quad \text{for all } F \subseteq E_K \\ & x(F) = p \\ & x_{jk} \in \{0,1\} \quad \text{for all } j \in N, k \in K \end{aligned}$$

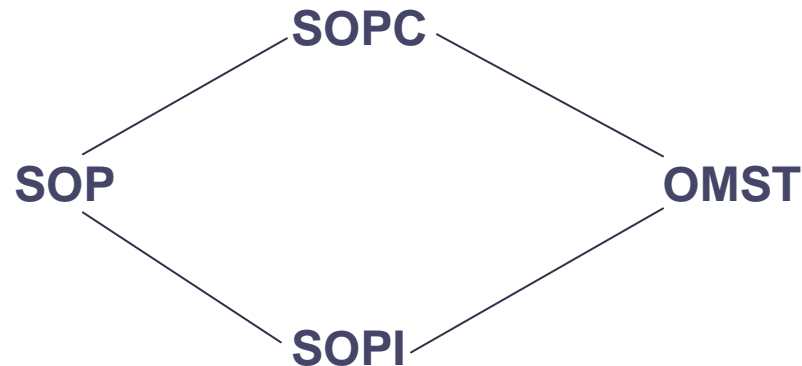
The OMSTP is a SOPIC with $d_j^k = \lambda_k c_j$ and $J = (E, \mathcal{H})$ given by forests in G .

Proposition:

The greedy algorithm yields an optimal solution to the OMSTP

Concluding remarks

- Discrete optimization problems with ordering



- The paper can be found at <http://www-eio.upc.es/~elena/doo.pdf>

Future work

- Characterization of P_{SOPI}
- Study other (easy) problems with order
- Study other (less easy) problems with order
- ...

Thank you for your attention!